

Arc-Disjoint Paths and Trees in 2-Regular Digraphs

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Abstract

An out-(in-)branching $B_s^+(B_s^-)$ rooted at s in a digraph D is a connected spanning subdigraph of D in which every vertex $x \neq s$ has precisely one arc entering (leaving) it and s has no arcs entering (leaving) it. We settle the complexity of the following two problems:

- Given a 2-regular digraph D , decide if it contains two arc-disjoint branchings B_u^+, B_v^- .
- Given a 2-regular digraph D , decide if it contains an out-branching B_u^+ such that D remains connected after removing the arcs of B_u^+ .

Both problems are NP-complete for general digraphs [1, 5]. We prove that the first problem remains NP-complete for 2-regular digraphs, whereas the second problem turns out to be polynomial when we do not prescribe the root in advance. We also prove that,

for 2-regular digraphs, the latter problem is in fact equivalent to deciding if D contains two arc-disjoint out-branchings. We generalize this result to k -regular digraphs where we want to find a number of pairwise arc-disjoint spanning trees and out-branchings such that there are k in total, again without prescribing any roots.

Keywords: Spanning tree, Out-branching, Mixed problem, Polynomial time, NP-Complete problem, 2-regular digraph.

1. Introduction

Every digraph will be finite. Notation will follow [2] unless stated otherwise. We recall the most relevant concepts below:

For a given digraph $D = (V, A)$ the **in-degree** $d^-(X)$ (**out-degree** $d^+(X)$) of the vertex set $X \subset V$ is the number of arcs entering (leaving) X . A digraph is k -**arc-strong** if $d^+(X) \geq k$ for every non-empty proper subset X of V .

When we **split** a vertex $v \in V$ into its ingoing part v^- and outgoing part v^+ we replace v by two new vertices v^- and v^+ and replace every arc $uv \in A$ ($vw \in A$) by the arc uv^- (v^+w).

A digraph $D = (V, A)$ is **k -regular** if every vertex $v \in V$ has out-degree $d^+(v) = k$ and in-degree $d^-(v) = k$.

A **Hamiltonian path (cycle)** in D is a directed path (cycle) that contains all vertices in V .

An **out-(in-)branching** $B_s^+(B_s^-)$ rooted at the vertex s in D is a connected spanning subdigraph of D in which every vertex $x \neq s$ has precisely one arc entering (leaving) it and s has no arcs entering (leaving) it.

A **spanning tree** in D is a spanning tree in the underlying graph $UG(D)$, that is the graph that appears when we disregard the direction of all arcs in D , turning them into edges. Note that $UG(D)$ may contain parallel edges. We say that a digraph D is **connected** when $UG(D)$ is a connected graph.

Given a digraph $D = (V, A)$, a vector $r : V \rightarrow \mathbb{N}_0$, that maps vertices to integers, is called a **root vector**

of D if there exists a set of arc-disjoint out-branchings in $D = (V, A)$ such that each vertex $v \in V$ is the root of exactly $r(v)$ of these out-branchings. We extend r to subsets of V by letting $r(X) = \sum_{v \in X} r(v)$ for all $X \subseteq V$. Note that $r(V)$ then denotes the total number of roots of out-branchings prescribed by r , assuming r is a root vector.

Graphs that contain k edge-disjoint spanning trees were characterized by Tutte

Theorem 1.1 (Tutte's Tree Packing Theorem). [12]
A graph $G = (V, E)$ contains k edge-disjoint spanning trees if and only if

$$e_{\mathcal{F}} \geq k(t - 1)$$

holds for every partition $\mathcal{F} = V_1, V_2, \dots, V_t$ of V , where $e_{\mathcal{F}}$ denotes the number of edges connecting different sets V_i, V_j .

A maximum collection of edge-disjoint spanning trees can be found in polynomial time by converting the problem to a matroid problem and then applying Edmonds' Matroid Partition Algorithm as explained in [7].

Digraphs that contain k arc-disjoint out-branchings have been characterized by Edmonds' Branching Theorem [8] from which one easily gets the following characterization of root-vectors:

Theorem 1.2. (T 2.12 in [6])

Let $D = (V, A)$ be a directed graph and $r : V \rightarrow \mathbb{N}$ a vector with $r(V) = k$.

Then r is a root vector if and only if

$$d_D^-(X) \geq k - r(X) \quad \text{for all non-empty } X \subseteq V.$$

Edmonds [8] gave a polynomial algorithm for finding these k arc-disjoint out-branchings given the roots specified by the root vector, while Frank [9] gave an algorithm for finding the maximum number of arc-disjoint out-branchings when we do not fix the roots.

In this paper we will study two related problems, both of which are NP-complete for general digraphs, and examine them on the restricted class of 2-regular digraphs.

Problem 1. Given a digraph D and vertices u, v (not necessarily distinct). Decide whether D has a pair of arc-disjoint branchings B_u^+, B_v^- .

Thomassen [11] conjectured in 1985 that for large enough integers r every r -arc-strong digraph contains arc-disjoint branchings B_v^+, B_v^- for every vertex v . This conjecture is wide open and it is only known that $r \geq 3$ must hold.

Thomassen also proved (see [1]) that Problem 1 is NP-complete for general digraphs. Thus it is of interest to study Problem 1 in special classes of digraphs. The first author proved that Problem 1 is polynomial for Tournaments [1]. We prove below that Problem 1 is NP-complete for 2-regular digraphs.

Inspired by the existence of good characterizations and algorithms (mentioned above) for respectively, the existence of k edge-disjoint spanning trees in a graph and the existence of k arc-disjoint out-branchings (with or without specified roots) in a digraph, Thomassé posed the following problem around 2006 (it appeared on the Hungarian problem page Egres open for several years).

Problem 2. Given a digraph D , decide whether it contains an out-branching B_u^+ such that D remains connected after removing the arcs of B_u^+ .

The first author and Yeo proved recently that Problem 2 is NP-complete for general digraphs [5]. In the meantime the idea of studying mixed problems where we are asking for (arc)-disjoint structures S, T in a digraph D where only S has to respect the orientation of arcs in D , gave inspiration for several papers, see e.g. [3, 4].

We prove that when we do not specify the root u in Problem 2 and D is 2-regular, then the problem becomes equivalent to deciding if the digraph contains 2 arc-disjoint out-branchings, hence making it polynomial. This contrasts the previous result that Problem 1 remained NP-complete even in the class of 2-regular digraphs. The complexity of the remaining case where we do specify the root is still open. In Section 5 we prove that a number of (seemingly) closely related problems are NP-complete even for 2-regular digraphs .

Packing two spanning subdigraphs in a 2-regular digraph will always require all but at most two of the arcs. So at first hand it seems that this restriction should make the problem tractable but as we mentioned above (and will prove in Sections 3, 4) this is only the case for Problem 2. In fact, under the assumption that $P \neq NP$, the following is a consequence of Problem 1 being NP-complete for 2-regular digraphs (see Theorem 3.1).

Corollary 1.3. *It is NP-complete to decide for a given digraph D on n vertices and $2n-2$ arcs, whether the arcs of D can be partitioned into an out branching B_u^+ and an in-branching B_v^- .*

In contrast to this, it follows from our discussion in the introduction that there is a polynomial algorithm for checking whether a given set of $2n-2$ edges in a graph on n vertices can be partitioned into two edge-disjoint spanning trees and there is also a polynomial algorithm for checking whether a set of $2n-2$ arcs in a digraph on n vertices can be partitioned into two arc-disjoint out-branchings with or without prescribed roots.

2. Hamiltonian Paths in 2-regular digraphs

We begin our investigation by considering a problem that requires all arcs, namely deciding if a 2-regular digraph contains two arc-disjoint Hamiltonian cycles. We point out that all of the results in this section have been proven by Yeo earlier but the proof of Theorem 2.2 was never published and since we use the technique illustrated below in our proof in Section 3 we have included it here for completeness. We also point out that Plesník [10] proved much earlier that both the Hamilton cycle and the Hamilton path problem are NP-complete already for planar 2-regular digraphs.

Theorem 2.1. [5, Theorem 6.1.3]

It is NP-complete to decide whether a given 2-arc-strong 2-regular digraph D contains two arc-disjoint Hamiltonian cycles.

Theorem 2.2. *It is NP-complete to decide whether a given 2-arc-strong 2-regular digraph contains two arc-disjoint Hamiltonian paths (with any number of specified end vertices).*

Proof. We will reduce from the problem of deciding if a 2-arc-strong 2-regular digraph contains two arc-disjoint Hamiltonian cycles, which according to Theorem 2.1 is NP-complete.

To do this we will use the Cycle Breaker Gadget shown in Figure 1. Notice that it is impossible to remove the arcs of a spanning (s, t) -path from the gadget without disconnecting some vertices.

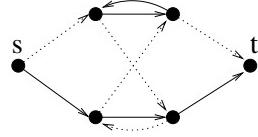


Figure 1: The Cycle Breaker Gadget. One of the two possible spanning (s, t) -paths is highlighted. Removing it disconnects 2 vertices.

Given a 2-arc-strong 2-regular digraph D construct D' by choosing a vertex $a \in D$ and splitting it into its in-going part a^- and its out-going part a^+ . Then construct D'' by adding a copy G of the Cycle Breaker Gadget to D' and identifying a^- with s and t with a^+ . Notice that this turns D'' into a 2-arc-strong 2-regular digraph.

To conclude the reduction we will argue that D'' contains two arc-disjoint Hamiltonian paths if and only if D contains two arc-disjoint Hamiltonian cycles.

For sufficiency assume that D contains arc-disjoint Hamiltonian cycles C and C' . Then C and C' can be considered as Hamiltonian (a^+, a^-) -paths in D' . Now adding the two arc-disjoint path fragments P and Q both covering the Cycle Breaker Gadget, see Figure 2, to C and C' considered in D'' gives an arc-disjoint pair of paths, where $C + P$ is a Hamiltonian (d, e) -path and $C' + Q$ is a Hamiltonian (b, c) -path of D' .

Conversely let P' and Q' be arc-disjoint Hamiltonian paths of D' . Neither can contain an (s, t) -path since that would disconnect part of the Cycle Breaker

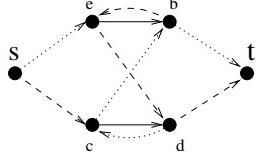


Figure 2: The path fragments P shown with dotted arcs and Q shown with dashed.

Gadget. But that implies that both must contain a (a^+, a^-) -path spanning all of D' translating into two arc-disjoint Hamiltonian cycles of D when a^+ and a^- are identified.

Since end vertices of the Hamiltonian paths are already forced to lie inside G , specifying one or more of them and requesting an arc-disjoint pair of a Hamiltonian (d, e) -path and a Hamiltonian (b, c) -path still gives the same reduction. \square

3. Arc-disjoint in- and out-branchings in k-regular digraphs

We now show that, using the Cycle Breaker Gadget, we can restrict the behavior of branchings in 2-regular digraphs. Given a 2-regular digraph D we immediately see that removing two arc-disjoint branchings would leave only two arcs in D . Suppose we specify two vertices $u, v \in D$ and assume arc-disjoint branchings B_u^+ and B_v^- exist.

Then the fact that both $d_{B_v^-}^+(x) = 1$ for all $x \in V - v$ and $d_{B_v^-}^+(v) = 0$ must hold implies that v is the only vertex that could have two out-going arcs in B_u^+ . Similarly only u can have two in-going arcs in B_v^- .

So our branchings may only really "branch" on u or v , in all other vertices they will behave just like paths.

Theorem 3.1. *Problem 1 (with or without fixed not necessarily distinct roots) is NP-complete for 2-arc-strong 2-regular digraphs.*

Proof. In order to prove NP-completeness we will reduce from the problem of deciding if a 2-arc-strong 2-regular digraph D contains two arc-disjoint Hamiltonian cycles.

Construct D' from D by choosing a vertex $a \in D$ and splitting it into its in-going part a^- and its out-going part a^+ . Then construct D'' by adding two copies G, G' of the Cycle Breaker Gadget to D' and identifying a^- with s, t with s' and t' with a^+ , see Figure 3.

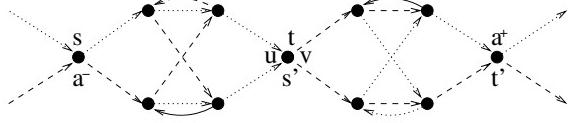


Figure 3: The in-branching fragment P shown with dotted arcs and the out-branching fragment Q shown with dashed.

Obviously there is no (s, t') -path covering all vertices of $G \cup G'$ which leaves a connected graph once the arcs of the path are removed. This implies that neither B_u^+ nor B_v^- can consist of a covering (s, t') -path when restricted to $G \cup G'$. By our observation above we know that our branchings can only behave differently than a path at the root of the other branching, so both u and v must be in $G \cup G'$. From this follows that, when restricted to arcs of D' , the branchings B_u^+ and B_v^- are arc-disjoint Hamiltonian (a^+, a^-) -paths in D' and thus arc-disjoint Hamiltonian cycles in D .

Conversely, if D contains two arc-disjoint Hamiltonian cycles C and C' , then those correspond to arc-disjoint (a^+, a^-) -paths in D' and adding branching fragments P and Q as in Figure 3 we can construct arc-disjoint in- and out-branchings in G' with the same root.

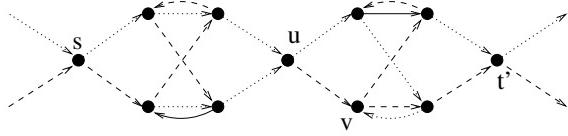


Figure 4: The in-branching fragment P shown with dotted arcs and the out-branching fragment Q shown with dashed.

If we add branching fragments P and Q as in figure 4 instead we get distinct roots. \square

Since 2-regular digraphs are also Eulerian this directly implies:

Corollary 3.2. *Problem 1 is NP-complete for Eulerian digraphs.*

With a bit of work we can extend our result further to k-regular digraphs.

Corollary 3.3. *Problem 1 is NP-complete for 2-arc-strong k-regular digraphs.*

Proof. In order to prove NP-completeness we will reduce from Problem 1 for 2-arc-strong 2-regular digraphs.

Let the digraph H have vertex set $\{b, c\}$ and arc set $\{bc, bc, cb, cb\}$.

Given a 2-arc-strong 2-regular digraph D construct D' by doing the following for every vertex $a \in V$. Split a into a^+ and a^- , add two copies of the arc a^-a^+ , add a copy of H and finally add $k-2$ copies of the arcs a^-b , ba^+ , a^+c and ca^- , see figure 5. Notice that D' is now a 2-arc-strong k-regular digraph.

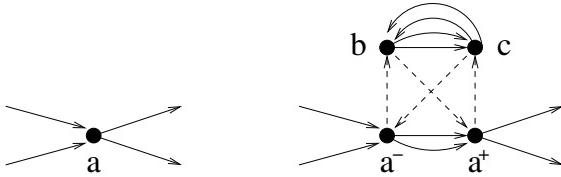


Figure 5: Construction of a k-regular digraph from a 2-regular one. The dashed arcs represent $k-2$ parallel arcs.

An arc-disjoint pair of an out-branching B_u^+ and an in-branching B_v^- in D can obviously be extended to an in- and out-branching pair in D' .

Conversely given arc-disjoint branchings B_u^+ and B_v^- in D' , simply undoing the steps we followed to construct D' from D will give a pair of branchings in D . \square

4. Arc-disjoint out-branchings and spanning trees

Since Problem 1 remained NP-complete for 2-regular digraphs one might wonder whether this is still the case for the mixed version, where the in-branching is replaced by a spanning tree in $UG(D)$.

We will show that Problem 2 is polynomially decidable for 2-regular digraphs, provided that the vertex u is not specified in advance. In fact we will consider a more general problem in k-regular digraphs.

Problem 3. Given a k -regular digraph D and $0 < l < k$ decide whether $UG(D)$ contains a collection of k -edge-disjoint spanning trees T_1, \dots, T_k , such that T_1, \dots, T_l are out-branchings in D (whereas we do not demand this from the remaining trees).

For $l = 0$ respectively $l = k$ we would get the purely undirected respectively purely directed problem characterized in Theorem 1.1 and Theorem 1.2 respectively, where we referenced polynomial algorithms that solve both pure problems. The simplest mixed problem (with $k = 2$ and $l = 1$) on the other hand is already NP-complete for general digraphs [5]. But for k-regular digraphs we have:

Theorem 4.1. *Let $D = (V, A)$ be a k -regular digraph then D contains k arc-disjoint out-branchings (with no restrictions on the roots) if and only if it contains k arc-disjoint spanning-trees.*

Proof. Since every out-branching is also a spanning tree necessity is obvious.

For sufficiency assume that D contains the k spanning trees T_1, \dots, T_k . Since $|A(T_i)| = |V| - 1$ we observe that the graph $D' = D \langle \bigcup_{i=1}^k T_i \rangle$, that is the union of the spanning trees, has exactly k arcs fewer than D . We denote these arcs by u_1v_1, \dots, u_kv_k so that $D' = D - \{u_1v_1, \dots, u_kv_k\}$.

Let $r(X) = \sum_{i=1}^k |X \cap \{v_i\}|$ for $X \subseteq V$ and notice that now every vertex v in D' has $d_{D'}^-(v) = k - r(v)$. We claim that r is a root vector of D' , which would immediately imply the existence of k arc-disjoint out-branchings since $r(V) = k$. To prove r is a root vector of D' it is sufficient, by Theorem 1.2, to prove that $d_{D'}^-(X) \geq k - r(X)$ holds for all $X \subset V$.

Since D' is the union of k spanning trees we have $|A_{D'}(X)| \leq k|X| - k$ for all $X \subseteq V$. On the other hand we know the in-degrees in D' so we can give the exact number of arcs as $|A_{D'}(X)| = k|X| - d_{D'}^-(X) - r(X)$.

Combining this we get

$$k|X| - d_{D'}^-(X) - r(X) = |A_{D'}(X)| \leq k|X| - k \\ \Rightarrow k - r(X) \leq d_{D'}^-(X)$$

So r is a root vector and since $r(V) = k$ Theorem 1.2 gives that there exist k arc-disjoint out-branchings rooted in D' . \square

This result implies that every k -regular digraph that contains a solution to Problem 3 regardless of l (where no roots are fixed) also contains k arc-disjoint out-branchings, since the solution to Problem 3 contains k arc-disjoint spanning trees. So we can decide the problem by either employing Frank's [9] algorithm for finding k arc-disjoint out-branchings without prescribed roots, or we employ Edmonds' algorithm [7] to find k arc-disjoint spanning trees, locate the roots as in the proof of Theorem 3 and then use Edmonds' branching algorithm [8] to find k arc-disjoint out-branching with these roots.

Corollary 4.2. *Problem 3 is polynomially solvable.*

One might be tempted to assume that regularity could be replaced by Eulericity in Theorem 4.1, but the digraph in Figure 6 shows that the theorem does not hold for Eulerian digraphs.

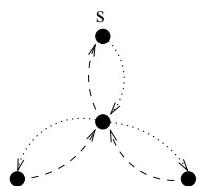


Figure 6: The out-branching B_s^+ shown with dotted arcs and the spanning tree T shown with dashed. Obviously it is not possible to find 2 arc-disjoint out-branchings in this digraph.

If we try to fix the roots, the k arc-disjoint out-branchings provided by Theorem 4.1 no longer necessarily constitute a solution as witnessed by the example in Figure 7. We can find two arc-disjoint out-branchings rooted at u and v respectively, but the indicated partition makes it obvious that the digraph allows for no out-branching rooted at s that leaves a connected subdigraph after removal.

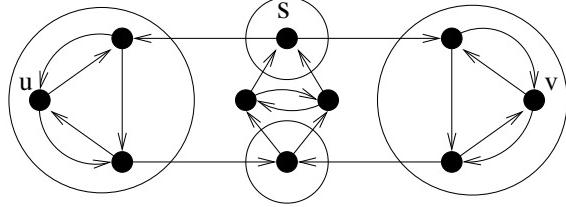


Figure 7: Every out-branching rooted at s will always disconnect this digraph upon removal.

5. Related NP-complete problems for 2-regular digraphs

While we saw above that there is a polynomial algorithm for deciding, in a 2-regular digraph D , the existence of an out-branching B_s^+ (whose root is not specified) such that $D - A(B_s^+)$ is connected, we could not resolve the question when we insist that s is specified. In what follows we will show that problems which are closely related are indeed NP-complete for 2-regular digraphs.

Let F be the digraph on 3 vertices $\{\alpha, \beta, \gamma\}$ and arcs $\{\alpha\beta, \beta\gamma, \gamma\alpha, \alpha\gamma, \gamma\beta\}$ (so F is obtained from the complete digraph on 3 vertices $\{\alpha, \beta, \gamma\}$ by deleting the arc $\beta\alpha$).

Theorem 5.1. *The following problems are NP-complete for 2-regular digraphs D*

- (P1) Given distinct vertices s, t ; does D have an (s, t) -path P such that $D - A(P)$ is connected?
 - (P2) Given distinct vertices s, t ; does D have an (s, t) -path P such that $D - A(P)$ is strongly connected?
 - (P3) Given distinct vertices s, t ; does D have an (s, t) -path P such that $D - A(P)$ contains an out-branching rooted at s ?

Proof:

We will show how to reduce 3-SAT to each of the problems (P1)-(P3) in polynomial time. We will use almost the same reduction for all three problems. The reduction is similar to one used in [5].

Let $W[u, v, p, q]$ be the digraph (the variable gadget) with vertices $\{u, v, y_1, y_2, \dots, y_p, z_1, z_2, \dots, z_q\}$ and the

arcs of the two (directed) (u, v) -paths $uy_1y_2 \dots y_p v$, $uz_1z_2 \dots z_q v$.

Let \mathcal{F} be an instance of 3-SAT with variables x_1, x_2, \dots, x_n and clauses C_1, C_2, \dots, C_m . We may assume that each variable x occurs at least once either in the negated form or non-negated in \mathcal{F} . The ordering of the clauses C_1, C_2, \dots, C_m induces an ordering of the occurrences of a variable x and its negation \bar{x} in these. With each variable x_i we associate a copy of $W[u_i, v_i, p_i, q_i]$ where x_i occurs p_i times and \bar{x}_i occurs q_i times in the clauses of \mathcal{F} . Identify end vertices of these graphs by setting $v_i = u_{i+1}$ for $i = 1, 2, \dots, n - 1$. Let $s = u_1$ and $t = v_n$.

Let D' be the digraph obtained in this way.

For each $i = 1, 2, \dots, m$ we associate the clause C_i with three of the vertices $V_i = \{a_{i,1}, a_{i,2}, a_{i,3}\}$ from the digraph D' above as follows:

Assume C_i contains variables x_j, x_k, x_l (negated or not). If x_j is not negated in C_i and this is the r 'th copy of x_j (in the order of the clauses that use x_j), then we identify $a_{i,1}$ with $y_{j,r}$ and if C_i contains \bar{x}_j and this is the g 'th occurrence of \bar{x}_j , then we identify $a_{i,1}$ with $z_{j,g}$. We make similar identifications for $a_{i,2}, a_{i,3}$. Thus D' contains all the vertices $a_{j,i}$ for $j \in [m], i \in [3]$.

The following Claim was proven in [5] but we include the easy proof for completeness.

Claim. D' contains an (s, t) -path P which avoids at least one vertex from $\{a_{j,1}, a_{j,2}, a_{j,3}\}$ for each $j \in [m]$ if and only if \mathcal{F} is satisfiable.

Proof of the Claim:

Suppose P is an (s, t) -path that avoids at least one vertex from $\{a_{j,1}, a_{j,2}, a_{j,3}\}$ for each $j \in [m]$. By construction, for each variable x_i , P traverses either the subpath $u_i y_{i,1} y_{i,2} \dots y_{i,p_i} v_i$ or the subpath $u_i z_{i,1} z_{i,2} \dots z_{i,q_i} v_i$. Now define a truth assignment by setting x_i false precisely when the first subpath is traversed for x_i . This is a satisfying truth assignment for \mathcal{F} since for any clause C_j at least one literal is avoided by P and hence becomes true by the assignment (the literals traversed become false and those not traversed become true). Conversely, given

a truth assignment for \mathcal{F} we can form P by routing it through all the false literals in the chain of variable gadgets. This proves the claim. \diamond

Let D be the 2-regular digraph we obtain from D' as follows:

For each clause C_i (to which we associated the vertices $V_i = \{a_{i,1}, a_{i,2}, a_{i,3}\}$ from the digraph D' above) we add 3 copies $F_{i,1}, F_{i,2}, F_{i,3}$ of F , with the vertices of the h th copy denoted by $\{\alpha_{i,h}, \beta_{i,h}, \gamma_{i,h}\}$ for $h = 1, 2, 3$. and the following arcs $\{a_{i,1}\alpha_{i,1}, \beta_{i,1}a_{i,2}, a_{i,2}\alpha_{i,2}, \beta_{i,2}a_{i,3}, a_{i,3}\alpha_{i,3}, \beta_{i,3}a_{i,1}\}$. Finally take two further copies F_1, F_2 with vertices $\{\alpha_1, \beta_1, \gamma_1\}$ and $\{\alpha_2, \beta_2, \gamma_2\}$ and add the arcs $\{t\alpha_1, t\alpha_2, \beta_1 s, \beta_2 s\}$.

Now it is easy to see that D has an (s, t) -path P such that removing the arcs of P leaves a connected digraph if and only if D' contains an (s, t) -path which avoids at least one vertex from each of the sets V_i , $i = 1, 2, \dots, m$ (P cannot enter any copy F' of F since that would disconnect these vertices from the rest after removing the arcs of P). Now the claim implies that (P1) is NP-complete.

To prove that (P2) is NP-complete we just need to show that if $D - A(P)$ is connected for some (s, t) -path P , then it is also strongly connected.

This follows from the construction since each copy of F is a strongly connected subgraph (on 3 vertices) and since at least one vertex of V_i , $i = 1, 2, \dots, m$ is left untouched by P it follows that $D - A(P)$ contains a cycle C through s, t formed by the arcs of $A(D') - A(P)$ and the path $t\alpha_1\gamma_1\beta_1s$ such that C contains at least one vertex from V_i for $i = 1, 2, \dots, m$. Clearly, as P does not enter any copy of F , we can attach F_2 and all of the copies $F_{i,j}$, $i \in [m], j \in [3]$ to this structure and obtain a strong spanning subdigraph of D which is arc-disjoint from P .

Finally, to prove that (P3) is NP-complete, we just need to observe that if P is any (s, t) -path in $UG(D)$ such that $D - A(P)$ is connected then P must be a directed (s, t) -path in D' and then, as we saw above, $D - A(P)$ is strongly connected and thus contains an out-branching from s . Clearly, if $D - A(P)$ contains an out-branching B_s^+ , then the (s, t) -path contained in B_s^+ must meet each of the sets V_i for

$i = 1, 2, \dots, m$, implying that P avoids at least one vertex from each V_i for $i = 1, 2, \dots, m$ so, by the claim, D is a yes instance for (P3) if and only if \mathcal{F} is satisfiable. \square

Notice that the only viable undirected (s, t) -path in the proof above is still a directed (s, t) -path in D' , so the problem is also NPC for undirected (s, t) -paths. Similarly if we delete F_1 from the graph and insert the arc ts instead we get a proof that the problems are NPC for directed cycles containing s instead of (s, t) -paths.

6. Concluding remarks

We did not completely settle Problem 2 for 2-regular digraphs since we had to leave open the case when the vertex u is fixed in advance:

Problem 4. What is the complexity of Problem 2 when D is 2-regular and the vertex u is part of the input?

The proof in [5] that Problem 2 is NP-complete for general digraphs involves constructing a digraph which is not the union of two arc-disjoint spanning trees. It seems difficult to modify that proof so that the digraph used is the union of two arc-disjoint spanning trees.

Problem 5. What is the complexity of Problem 2 when the input digraph $D = (V, A)$ is the union of two arc-disjoint spanning trees?

Note that, just as Corollary 1.3 followed directly from Theorem 3.1, we again have that if Problem 4 is NP-complete, then so is Problem 5 but it may be the case that Problem 5 is NP-complete while Problem 4 is still polynomially solvable.

The following result from [5] indicates that Problem 4 could be NP-complete. Note that here the digraph H has either n or $n + 1$ arcs since we need $n - 1$ for the connected remainder.

Theorem 6.1. [5]

It is NP-complete to decide whether a strongly connected 2-regular digraph D contains a spanning strong subdigraph H so that $D - A(H)$ is connected.

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